

S P E C I F I C A T I O N

Pseudo Noise Generator

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SPECIFICATION  
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## FIELD OF THE INVENTION

The present invention relates to a pseudo noise generator employed to generate pseudo noise, etc. for evaluate immunity of electric devices against electro-magnetic interference waves.

## DESCRIPTION OF PRIOR ART

In case of combining many kinds of noise, or of thermal noise or city noise to one another, amplitude of noise has the Gaussian distribution. In order to simulate generated noises to the Gaussian distribution, there have been proposed a noise generator for generating white-Gaussian noise by the use of a noise diode.

In narrow-band digital communication systems, there is a correlation between a bit error rate of the narrow-band digital communication system and an Amplitude Probability Distribution (referred to A P D) of electro-magnetic interference waves. Moreover, a report was published that the bit error rate of the communication system can be evaluated from the A P D of electro-magnetic interference waves. In this report, noise determined by a specified A P D is generated by the use of an arbitrary distribution random number generator (See: Paper Journal of The Institute of Electrical Communication Engineers of Japan (A), vol. J 70-A, No.11, ppl681 -1690, Nov.1987 ).

Important parameters for defining characteristic nature of noise are the A P D , a Crossing Rate Distribution (referred to C R D ), a Pulse Duration Distribution (referred to P D D ), and a Pulse Spacing Distribution (referred to P S D ) etc. These parameters will now be described with reference to Fig.13.

The A P D is defined as a time rate where the instantaneous value of a signal, such as electro-magnetic interference waves, exceeds a predetermined value, to show a total time length of the instantaneous value exceeding a level  $E_k$  in a test time period of  $T_0$  . The C R D

is defined as a number of crossings per a unit time where the instantaneous value of the signal crosses the specified level  $E_k$  to a positive direction ( or a negative direction ).

The P D D is defined as a probability distribution of a time  $W_1$  (k) where the instantaneous value of the signal exceeds a level  $E_k$  in a test time period of  $T_0$ . On the contrary, the P S D is defined as a probability distribution of a time  $Z_1$  (k) where the instantaneous value of the signal lowers a level  $E_k$  in a test time period of  $T_0$ . In other words, the P D D and the P S D are probability distributions of time lengths from a time where the instantaneous value of the signal crosses the threshold level to a just succeeding time where the instantaneous value of the signal crosses the threshold level.

Moreover, a Probability Density Function is defined as a distribution of the level  $E_k$  in the test time period of  $T_0$ .

In a pseudo noise generator, the dispersion and the average of noise can be specified, but the A P D of noise cannot be specified since distribution of noise is limited to the white-Gaussian noise.

In an arbitrary distribution random generator, the A P D of noise be specified to generate noise with an arbitrary A P D. However, noise from the arbitrary distribution random generator assumes independent events having no time-correlation. On the contrary, noise from electronic Ranges and ordinary electronic devices assumes non-independent events dependent to the period of a source voltage and the period of timing clock pulses. Accordingly, the C R D, the P D D and the P S D of noise from the arbitrary distribution random generator are different from the C R D, the P D D and the P S D of noise of non-independent events having time-correlation.

## SUMMARY OF THE INVENTION

An object of the present invention is to provide a pseudo noise

generator capable of specifying a Pulse Duration Distribution and a Pulse Spacing Distribution of noise at a specified amplitude level K in addition to an Amplitude Probability Distribution of noise.

5 To this end, a pseudo noise generator of the present invention comprises :

a first arbitrary random number generator for generating two groups of first random number signals respectively corresponding to divided Amplitude Probability Distributions, which are obtained by dividing a specified Amplitude Probability Distribution into two parts at a  
10 specified level;

a second arbitrary random number generator for generating two groups of second random number signals respectively defined by a specified Pulse Duration Distribution and a specified Pulse Spacing Distribution at the specified level;

15 control means for selecting ones of said two groups of first random number signals in accordance with said specified Pulse Duration Distribution and said specified Pulse Spacing Distribution defined at the specified level; and

a D/A converter for converting the selected signals to pseudo noise  
20 of analog value; said pseudo noise being generated in accordance with said Amplitude Probability Distribution, and said specified Pulse Duration Distribution and said specified Pulse Spacing Distribution at the specified level.

#### BRIEF DESCRIPTION OF THE DRAWINGS

25 The present invention will be described in detail below with reference to the accompanying drawings, in which;

Fig.1 is a block diagram explanatory of the principle of the present invention;

Fig.2 is a block diagram illustrating an embodiment of the present

invention;

Fig.3 is a block diagram illustrating an example of an arbitrary distribution random number generator 1;

Fig.4 is a block diagram illustrating an example of each bit generator employed in arbitrary distribution random number generators 1 and 2;

Fig.5 is a graph showing arrangement of data in a memory in the bitgenerator;

Fig.6 is a block diagram illustrating an example of the arbitrary distribution random number generator 1 and time charts explanatory of operations of the same;

Fig.7 is a block diagram illustrating an example of the arbitrary distribution random number generator 2 and time charts explanatory of operations of the same;

Fig.8 is a block diagram illustrating an example of the controller 3 and time charts explanatory of operations of the same;

Fig.9 shows characteristic curves illustrating conventional test results of the Amplitude Probability Distribution (a) , the Crossing Rate Distribution (b) and the Probability Density Function (c) of electro-magnetic interference waves from electronic Ranges and pseudo noise in case of specifying the electro-magnetic interference waves from electronic Ranges and the Amplitude Probability Distribution;

Fig.10 shows characteristic curves illustrating conventional test results of the Pulse Duration Distribution of electro-magnetic interference waves from electronic Ranges and pseudo noise in case of specifying the electro-magnetic interference waves from electronic Ranges and the Amplitude Probability Distribution;

Fig.11 shows characteristic curves illustrating the present invention's test results of the Amplitude Probability Distribution (a) ,

the Crossing Rate Distribution (b) and the Probability Density Function (c) of electro-magnetic interference waves from electronic Ranges and pseudo noise in case of specifying the electro-magnetic interference waves from electronic Ranges and the Amplitude Probability Distribution;

Fig.12 shows characteristic curves illustrating the present invention's test results of the Pulse Duration Distribution (a) and the Pulse Spacing Distribution (b) of electro-magnetic interference waves from electronic Ranges and pseudo noise in case of specifying the electro-magnetic interference waves from electronic Ranges and the Amplitude Probability Distribution; and

Fig.13 shows time charts explanatory of technical terms employed in this specification.

#### DETAILED DESCRIPTION

With reference to Fig.1, the principle of the pseudo noise generator of the present invention provided in accordance with a specified A P D, a specified P D D and a specified P S D will first be descrined. This pseudo noise generator of the present invention comprises four arbitrary distribution random number generators APD<sub>1</sub>, APD<sub>2</sub>, PDD<sub>0</sub> and PSD<sub>0</sub> and a selector SL.

In this case, setting of the P D D and the P S D is one point of the same level k; and an arbitrary distribution random number generator APD<sub>1</sub> for generating random numbers of values exceeding the level k and an arbitrary distribution random number generator APD<sub>2</sub> for generating random numbers of values under the level k are provided. These arbitrary distribution random number generators APD<sub>1</sub> and APD<sub>2</sub> are switched to meet with the specified P D D and P S D to generate random number codes meeting with required A P D, P D D and P S D. Namely, generation of pseudo noise according to the present invention is

performed as described below:

① Binary codes  $i_1$  of  $N$  bits included in a pulse duration distribution  $pdd(i_1)$  is generated from the arbitrary distribution random number generator  $PDD_0$  to determine a pulse duration length  $T_{11}$ .  
5 During this pulse duration length  $T_{11}$ , binary codes  $x_1$  of  $M$  bits included in an amplitude probability distribution  $apd_1(x_1)$  is generated from the arbitrary distribution random number generator  $APD_1$  to obtain pseudo noise  $x$ .

② After the end of the pulse duration  $T_{11}$ , binary codes  $i_2$  of  
10  $N$  bits included in a pulse duration distribution  $psd(i_2)$  is generated from the arbitrary distribution random number generator  $PSD_0$  to determine a pulse duration length  $T_{12}$ . During this pulse duration length  $T_{12}$ , binary codes  $x_2$  of  $M$  bits included in an amplitude probability distribution  $apd_2(x_2)$  is generated from the arbitrary  
15 distribution random number generator  $APD_2$  to obtain an output of pseudo noise  $x$ .

③ The above steps ① and ② are alternately performed.

In accordance with the above operations, binary numbers  $x$  of  $M$  bits meeting with the required A P D in addition to the required P D D  
20 and P S D are generated, and digital-to-analog converted to obtain an output of the pseudo noise generator.

As mentioned above, in the pseudo noise generator of the present invention, an arbitrary distribution random number generator 100 for generating exclusively the binary codes  $x_1$  of  $M$  bits included in an  
25 amplitude probability distribution  $apd_1(x_1)$  or the binary codes  $x_2$  of  $M$  bits included in an amplitude probability distribution  $apd_2(x_2)$  is employed in place of an arbitrary distribution random number generator for generating a signal  $x$  included in a specified amplitude probability distribution ( $x$ ).

A controller for outputting a memory selecting signal s is provided to select a first state generating the binary codes  $x_1$  or a second state generating the binary codes  $x_2$ .

5 An arbitrary distribution random number generator 200 is provided for generating exclusively the binary codes  $i_1$  of M bits included in a pulse duration distribution  $pdd_1(i_1)$  or the binary codes  $i_2$  of M bits included in a pulse spacing distribution  $psd(i_2)$  to generate the memory selection signal s employed in the controller. Setting of the  $pdd_1(i_1)$  and the  $psd(i_2)$  is one point of the same level k.

10 A digital-analog converter is provided to convert the binary codes x ( $x_1$  or  $x_2$ ) to an analog value. An analog signal converted in the digital-analog converter is applied to a communication system through a cable or is radiated from an antenna by shifting the frequency band of the thereof by the use of an up-converter.

15 In the pseudo noise generator according to the present invention, the binary codes  $i_1$  included in a pulse duration distribution  $pdd_1(i_1)$  is generated from the arbitrary distribution random number generator 200. During the pulse duration length  $T_{11}$  corresponding to the binary codes  $i_1$ , the binary codes  $x_1$  included in an amplitude probability distribution  $apd_1(x_1)$  is generated from the arbitrary distribution random number generator 100. Thereafter, the binary codes  $i_2$  included in a pulse spacing distribution  $psd(i_2)$  is generated from an arbitrary distribution random number generator 200. During the pulse duration length  $T_{12}$ , the binary codes  $x_2$  included in an amplitude probability distribution  $apd_2(x_2)$  is generated from the arbitrary distribution random number generator 100. Switching between the amplitude probability distribution  $apd_1(x_1)$  and the amplitude probability distribution  $apd_2(x_2)$  and switching between the pulse duration distribution  $pdd_1(i_1)$  and the pulse spacing distribution  $psd(i_2)$  are

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performed in accordance with the memory selection signal s from the selector SL.

Since the amplitude level of the binary codes  $x_1$  included in the amplitude probability distribution  $apd_1(x_1)$  exceeds the value k while  
5 the amplitude level of the binary codes  $x_2$  included in the amplitude probability distribution  $apd_2(x_2)$  lowers the value k, the setting of the P D D and the P S D at the amplitude level k is carried out in the pulse duration length  $T_{11}$  and the pulse duration time length  $T_{12}$ , respectively.

10 The amplitude probability distribution  $apd_1(x_1)$  and the amplitude probability distribution  $apd_2(x_2)$  are calculated from the amplitude probability distribution  $apd(x)$ . If the amplitude probability distribution  $apd(k)$ , the pulse duration distribution  $pdd_1(i_1)$  and the pulse spacing distribution  $psd(i_2)$  meet with the condition defined by  
15 an equation (1), the binary codes x ( $x_1$  or  $x_2$ ) generated from the arbitrary distribution random number generator 100 is included in the amplitude probability distribution  $apd(x)$ . In this case, the notation M is a number of bits of the binary codes x ( $x_1$  or  $x_2$ ) generated from the arbitrary distribution random number generator 100,  
20 and the notation N is a number of bits of the binary codes i ( $i_1$  or  $i_2$ ).

$$\sum_{i_2=0}^{2^N-1} psd(i_2) T_{12} = \frac{1 - apd(k)}{apd(k)} \sum_{i_1=0}^{2^N-1} psd(i_1) T_{11} \quad (1)$$

In the present invention, designation of the pulse duration  
25 distribution P D D and the pulse spacing distribution P S D is not limited to designation of distribution having a distribution duration, but can be included designation of a special example of distribution having a defined value, as far as meeting with the condition of the equation (1).

(Embodiments)

In Fig.2, an embodiment of the pseudo noise generator of the present invention is illustrated. In this embodiment, arbitrary distribution random number generators of eight bits are employed. The embodiment comprises an arbitrary distribution random number generator 1, an arbitrary distribution random number generator 2, a controller 3 and a digital-to-analog (D/A) converter 4. The arbitrary distribution random number generator 1 is controlled by clock (1) and the memory selection signal s from the controller 3. The arbitrary distribution random number generator 2 is controlled by clock (2) and the memory selection signal s which are applied from the controller 3. The D/A converter 4 converts the binary codes x to an analog signal.

With reference to Fig.3, the arbitrary distribution random number generator 1 comprises eight bit generators 1-1 to 1-8 and eight latch circuits 1-11 to 1-18 to generate the binary codes x ( $x_1$  or  $x_2$ ) of eight bits, in which the eight bit generators 1-1 to 1-8 and eight latch circuits 1-11 to 1-18 are alternately connected in cascade so as to actuate each of the bit generators 1-1 to 1-8 at the rise-up instants of each clock pulses.

As shown in Fig.4, each of the eight bit generators 1-1 to 1-8 comprises a uniform random number generator 11, a memory 12, and a comparator 13. In the memory 12, data y employed for determining each bit ( $a, b, \dots, h$ ) of the binary codes x ( $x_1$  or  $x_2$ ) are stored as follows.

The amplitude probability distribution  $apd_1(x_1)$  corresponds to a limited part of the binary codes x ( $x_1$ ) in the amplitude probability distribution  $apd(x)$  to generate the binary code  $x_1$ , which is included in a range  $k \leq x$ . The amplitude probability distribution  $apd_1(x_1)$  is defined in an equation (2).

$$\text{apd}_1(x_1) = \begin{cases} 1 & (0 \leq x_1 < k) \\ \frac{\text{apd}(x_1)}{\text{apd}(k)} & (k \leq x_1 \leq 2^8 - 1) \end{cases} \quad (2)$$

5 The amplitude probability distribution  $\text{apd}_2(x_2)$  corresponds to a limited part of the binary codes  $x_2$  in the amplitude probability distribution  $\text{apd}(x)$  to generate the binary code  $x_2$ , which is included in a range  $k > x$ . The amplitude probability distribution  $\text{apd}_2(x_2)$  is defined in an equation (3).

$$\text{apd}_2(x_2) = \begin{cases} \frac{\text{apd}(x_2) - \text{apd}(k)}{\text{apd}(0) - \text{apd}(k)} & (0 \leq x_2 < k) \\ 0 & (k \leq x_2 \leq 2^8 - 1) \end{cases} \quad (3)$$

The amplitude probability distribution  $\text{apd}_1(x_1)$  and the amplitude probability distribution  $\text{apd}_2(x_2)$  calculated in accordance with the steps defined in the equations (2) and (3) are converted to Conditional Probabilities  $\text{pc}_1(j, r)$ ,  $\text{pc}_2(j, r)$  and then stored as data\_y for determining each bit in the memory 5 as shown in Fig.5. In this case, the notation  $j = 1, 2, \dots, 8$  and the notation  $r = 0, 1, \dots, 2^{j-1} - 1$ .

$$\text{pc}_1(j, r) = \frac{\text{apd}_1((2r+1) \cdot 2^{M-j}) - \text{apd}_1((2r+2) \cdot 2^{M-j})}{\text{apd}_1((2r) \cdot 2^{M-j}) - \text{apd}_1((2r+2) \cdot 2^{M-j})} \quad (4)$$

$$\text{pc}_2(j, r) = \frac{\text{apd}_2((2r+1) \cdot 2^{M-j}) - \text{apd}_2((2r+2) \cdot 2^{M-j})}{\text{apd}_2((2r) \cdot 2^{M-j}) - \text{apd}_2((2r+2) \cdot 2^{M-j})} \quad (5)$$

With reference to Figs.6 and 4, operations of the arbitrary distribution random number generator 1 will be described.

In the arbitrary distribution random number generator 1, bit data

A( $s_1$ ,  $a_1$ ) from a first bit generator 1-1 is applied through a latch 1-11 to a second bit generator 1-2 at the rise instant of each pulse of the clock (1). The signal  $s_1$  is the memory selection signal s, and the signal  $a_1$  is a first bit a of the binary codes x from the first  
5 bit generator 1-1.

At the same time as the rise instant of each pulse of the clock (1), the uniform random number generator 11 in the second bit generator 1-2 having the construction shown in Fig.4 generates uniform random numbers z. Thereafter, the data y employed in the second bit  
10 generator 1-2 are read out from the memory 12 by the use of the bit data A( $s_1$ ,  $a_1$ ) as address data. The data y and z are compared with each other, so that a second bit  $b_1$  of an arbitrary distribution number x at the output of the comparator 12 assumes a state "1" in case of  $y < z$ , while the second bit  $b_1$  assumes the state "0" in case of  $y \geq z$ .

15 The second bit generator 1-2 applies bit data B ( $s_1$ ,  $a_1$ ,  $b_1$ ) to a third bit generator 1-3 (not shown) at the just succeeding rise instant of each pulse of the clock (1). New bit data of ( $s_2$ ,  $a_2$ ) are applied from the first bit generator 1-1 to the second bit generator 1-2, which generates second bit data  $b_2$  at the just succeeding clock pulse in  
20 response to the bit data of ( $s_2$ ,  $a_2$ ).

Each of other bit generators generates corresponding bit data from the bit data applied from a bit generator of the just preceding stage, so that the bit data applied from the bit generator of the just preceding stage are combined with bit data generated at the bit  
25 generator of instant stage, the combined bit data are applied to a bit generator of the just succeeding state.

In this case, since the first bit generator 1-1 is of a first stage, only the memory selection signal s is applied to the first bit generator 1-1 as shown in Fig.6 in place of bit data to be applied from a just

preceding stage. Since a bit generator 1-8 is of a last stage, the memory selection signal s is not necessarily applied to a next stage. Accordingly, data x ( $x = a, b, \dots, h$ ) from which the memory selection signal s is removed are applied to the D/A converter as shown in Fig.6.

5 In accordance with the operations mentioned above, the arbitrary distribution random number generator 1 generates the binary code  $x_1$  included in the amplitude probability distribution  $apd_1(x_1)$  in synchronism with the pulse of the clock (1) in case of the state "1" of the memory selection signal s from the controller 3. In case of the  
10 state "0" of the memory selection signal s from the controller 3, the binary code  $x_2$  included in the amplitude probability distribution  $apd_2(x_2)$  in synchronism with the pulse of the clock (1).

With reference to Fig.7, construction and operations of the arbitrary distribution random number generator 2 will now be described.  
15 This arbitrary distribution random number generator 2 generates, alternately, the binary codes  $i_1$  employed for determining the pulse duration time length  $T_{i_1}$  of the binary codes  $x_1$  included in the amplitude probability distribution  $apd_1(x_1)$ , or the binary codes  $i_2$  employed for determining the pulse duration length  $T_{i_2}$  of the binary  
20 codes  $x_2$  included in the amplitude probability distribution  $apd_2(x_2)$ .

This arbitrary distribution random number generator 2 has the construction and operations similar to those of the arbitrary distribution random number generator 1. In this arbitrary distribution random number generator 2, however, the clock (2) is  
25 employed in place of the clock (1). In each memory 12 of bit generators 2-1 to 2-8, data for generating the pulse duration distribution  $pdd_1(i_1)$  and the pulse spacing distribution  $psd(i_2)$  in place of the amplitude probability distribution  $apd_1(x_1)$  and the amplitude probability distribution  $apd_2(x_2)$ . From a function  $n(T_{i_1})$

) of a number  $n$  and the pulse duration length  $T_{11}$ , the pulse duration distribution  $pdd_1(i_1)$  can be calculated in accordance with an equation (6).

$$pdd(i_1) = \frac{n(T_{11})}{\sum n(T_{11})} \quad (6)$$

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The pulse spacing distribution  $psd(i_2)$  is calculated in the similar manner to those of the pulse duration distribution  $pdd_1(i_1)$  in accordance with an equation (7).

$$psd(i_2) = \frac{m(T_{12})}{\sum m(T_{12})} \quad (7)$$

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The pulse spacing distribution  $psd(i_2)$  and the pulse duration distribution  $pdd_1(i_1)$  are converted in accordance with equations (8) and (9) to trial conditional probability values  $pc_1(j,r)$  and  $pc_2(j,r)$ , which are stored in the memory of the arbitrary distribution random number generator 2. In this case, the binary codes  $i_1$  and  $i_2$  are defined by bit signals  $j$  and  $r$  in the equation (8). However, the bit signals  $j$  and  $r$  are defined as  $j=1,2,\dots,8$  and  $r=0,1,\dots,2^{j-1}-1$ . The bit signal  $r$  is determined by trial results until the  $(j-1)$ -th trial. Data arrangement in the memory 12 of each bit generation 2-1 to 2-8 is shown in Fig.5 in the similar manner to the arbitrary distribution random number generator 1.

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$$pc_1(j,r) = \frac{pdd((2r+1) \cdot 2^{8-j}) - pdd((2r+2) \cdot 2^{8-j})}{pdd(2r \cdot 2^{8-j}) - pdd((2r+2) \cdot 2^{8-j})} \quad (8)$$

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$$pc_2(j,k) = \frac{psd((2r+1) \cdot 2^{8-j}) - psd((2r+2) \cdot 2^{8-j})}{psd(2r \cdot 2^{8-j}) - psd((2r+2) \cdot 2^{8-j})} \quad (9)$$

With reference to Fig.8, the controller comprises an I-T converter

20 receiving the binary codes  $i$ , a down counter 21 of 32 bits receiving output data  $T_*$  of the I-T converter 20 under control of the clock (1), a memory selection signal generator 22 receiving the carry output  $c$  of the down counter 21, and a clock generator 23 receiving the carry output  $c$  of the down counter 21 and the clock (1). The I-T converter 20 is a memory for data of thirty-two bits which is controlled with address codes  $i$  ( $= i_1$  or  $i_2$ ) of eight address bits to read out the stored data  $T_*$  ( $*$  =  $i_1$  or  $i_2$ ).

The controller 3 sets the memory selection signal  $s$  to the state "1" during the time length  $T_{i1}$  and to the state "0" during the time length  $T_{i2}$  to control the pulse duration length  $T_{i1}$  of an arbitrary distribution random number included in the amplitude probability distribution  $apd_1(x_1)$  and the pulse duration length  $T_{i2}$  of an arbitrary distribution random number included in the amplitude probability distribution  $apd_2(x_2)$  as shown in time charts of Fig.8.

If the memory selection signal  $s$  assumes the state "1", the counting value  $c$  of the down counter 21 is decreased by "1". In this period of the state "1", the arbitrary distribution random number generator 1 generates the amplitude probability distribution  $apd_1(x_1)$ .

When the counting value  $c$  of the down counter 21 reaches the zero state, the memory selection signal  $s$  assumes the state "0". In this period of the state "0", the arbitrary distribution random number generator 1 generates the amplitude probability distribution  $apd_2(x_2)$ .

In response to the zero state of the down counter 21, the pulse duration length  $T_{i2}$  is newly set in the down counter 21 at the rise instant of a just succeeding pulse of the clock (1). At the same time as the zero state of the down counter 21, a pulse  $p_2$  of the clock (2) comes at the rise instant of a just succeeding pulse of the clock (1) to actuate the arbitrary distribution random number generator 2 so as

to newly generate the binary code  $i_1'$ , which is converted to the output data  $T_1'$  in the I-T converter 20.

The controller 3 decreases successively the counting value of the down counter 21 in the similar manner to the case of the state "1" of the memory selection signal  $\underline{s}$ . In this period, the arbitrary distribution random number generator 1 generates the binary code  $x_2$  included in the amplitude probability distribution  $apd_2(x_2)$ . When the counting value of the down counter 21 reaches the zero state, the state "0" of the memory selection signal  $\underline{s}$  is reversed to the state "1", so that the output data  $T_1'$  is set to the down counter 21 to be restored to the initial state.

As mentioned above, the pulse duration length  $T_{11}$  and the pulse duration length  $T_{12}$  are alternately set to the down counter 21 in the controller 3 to alternately reverse the state "1" and the state "0" of the memory selection signal  $\underline{s}$  in response to the pulse duration length  $T_{11}$  and the pulse duration length  $T_{12}$ , respectively, so that the binary code  $x_1$  included in the amplitude probability distribution  $apd_1(x_1)$  and the binary code  $x_2$  included in the amplitude probability distribution  $apd_2(x_2)$  are alternately generated.

Simulation results of pseudo noise generation by a computer associated with the device of the present invention will be described. In this simulation test, the A P D, the P D D and the P S D of the electro-magnetic interference waves from an electronic range are measured and adopted. Since the measured values of the A P D, the P D D and the P S D are obtained from actual electro-magnetic interference waves, these measured values meet necessarily with the condition of the equation (1). Figs. 9A, 9B, 9C and Fig. 10 are simulation results of pseudo noise distributions by conventional devices, where only the electro-magnetic interference waves from an electronic



range and the A P D are specified. A characteristic curve of heavy line in Fig.9A indicates the A P D of the electro-magnetic interference waves from an electronic range and is adopted as the A P D by conventional devices. A characteristic curve of heavy line in Fig.9B indicates the C R D of the electro-magnetic interference waves from an electronic range, and a characteristic curve of heavy line in Fig.9C indicates the P D F of the electro-magnetic interference waves from an electronic range. By using marks  $\bigcirc$ ,  $\triangle$  and  $\times$  in Figs.9A, 9B and 9C, the A P D, the C R D and the P D F of the pseudo noise are illustrated. In Fig.10, the P D D of the actual electro-magnetic interference waves and the P D D of the pseudo noise are illustrated with a heavy line and marks  $\bullet$ , respectively. As understood from Figs. 9A, 9B, 9C and 10, the actual electro-magnetic interference waves and the pseudou noise are substantially agreeable with each other with respect to the A P D and the P D F but different from each other with reference to the C R D and the P D D.

On the contrary, Figs. 11A, 11B and 11C and Fig.12 show test result of pseudou noise according to the present invention in case of specifying the actual electro-magnetic interference waves and the A P D, the P D D and the P S D. A heavy line in Fig.11A indicates the A P D of the actual electro-magnetic interference waves and is employed as a specified value of the A P D. A heavy line in Fig.11B indicates the C R D of the actual electro-magnetic interference waves, while a heavy line in Fig.11C indicates the P D F of the actual electro-magnetic interference waves. In Figs. 11A, 11B and 11C, marks  $\bigcirc$ ,  $\triangle$  and  $\times$  indicates the A P D, the C R D and the P D F of the actual electro-magnetic interference waves, respectively. In Fig.12A, a heavy line indicates the P D D of the actual electro-magnetic interference waves, while a heavy line indicates the P S D of the actual electro-magnetic

interference waves. In Figs.12A and 12B, the P D D and the P S D of pseudou noise generated in case of the P D D and the P S D specified are indicated by marks  $\bigcirc$  and  $\times$ , respectively. As understood from Fig. 11A to Fig.12B, the actual electro-magnetic interference waves and the pseudou noise are substantially agreeable with each other with respect to the P D D and the P S D in addition to the A P D and the P D F but still different from each other with reference to the C R D .

As mentioned above in details, noise of non-independent events having time-correlation of desired characteristics can be generated in accordance with the present invention of relatively simplified construction and control operations by designating the amplitude probability distribution  $apd(x)$ , the pulse duration distribution  $pdd_1(i_1)$  and the pulse spacing distribution  $psd(i_2)$ . Therefore, merits of the present invention are very effective for evaluate immunity of electric devices against electro-magnetic interference waves.